## Solutions for the AAC Competition Problems 2012

1. Let A be a square matrix of odd size. Prove that A is singular if and only if it can be carried to -A by elementary operations of adding a multiple of one row to another row.

**Solution:** The indicated elementary operations preserve the determinant, so a nonsingular matrix A cannot be carried to -A since  $\det(-A) = (-1)^n \det(A) = -\det(A) \neq \det(A)$ , where n is the size of A.

Now assume that A is singular. Then Gaussian elimination allows us to carry A to a matrix B that has a zero row at the bottom. Hence  $B = E_s \cdots E_1 A$  where  $E_k$ ,  $k = 1, \ldots, s$ , are elementary matrices corresponding to adding a multiple of one row to another row. It is easy to carry B to -B. Indeed, for each  $i = 1, \ldots, n-1$ , we can add row i to the last row, then subtract the double of the new last row from row i and finally add the new row i to the last row. Thus -B = UB where U is the product of 3(n-1) elementary matrices. It remains to observe that  $-A = E_1^{-1} \cdots E_s^{-1}(-B)$  and the matrices  $E_k^{-1}$  are also elementary. Hence

$$-A = E_1^{-1} \cdots E_s^{-1} U E_s \cdots E_1 A,$$

as required.

**2**. Let  $\xi$  be a primitive complex *n*-th root of unity,  $n \ge 2$ . Determine the *n*-th power of the following  $n \times n$  matrix:

(1	-1	0	0	0	 0	0 \	
0	ξ	$^{-1}$	0	0	 0	0	
0	0	$\xi^2$	-1	0	 0	0	
							•
0	0	0	0	0	 $\xi^{n-2}$	-1	
$\setminus -1$	0	0	0	0	 0	$\xi^{n-1}$	

**Solution**: Let us first compute the characteristic polynomial  $\chi(t)$  of the given matrix A:

$$\chi(t) = \det(A - tI) = \begin{vmatrix} 1 - t & -1 & 0 & 0 & 0 & \dots & 0 & 0 \\ 0 & \xi - t & -1 & 0 & 0 & \dots & 0 & 0 \\ 0 & 0 & \xi^2 - t & -1 & 0 & \dots & 0 & 0 \\ \dots & & & & & \\ 0 & 0 & 0 & 0 & 0 & \dots & \xi^{n-2} - t & -1 \\ -1 & 0 & 0 & 0 & 0 & \dots & 0 & \xi^{n-1} - t \end{vmatrix}.$$

Using expansion down the last column, we obtain:

$$\begin{split} \chi(t) = \begin{vmatrix} 1-t & -1 & 0 & 0 & 0 & \dots & 0 \\ 0 & \xi - t & -1 & 0 & 0 & \dots & 0 \\ 0 & 0 & \xi^2 - t & -1 & 0 & \dots & 0 \\ \dots & & & & & \\ 0 & 0 & 0 & 0 & 0 & \dots & -1 \\ -1 & 0 & 0 & 0 & 0 & \dots & 0 \end{vmatrix} \\ & + (\xi^{n-1} - t) \begin{vmatrix} 1-t & -1 & 0 & 0 & 0 & \dots & 0 \\ 0 & \xi - t & -1 & 0 & 0 & \dots & 0 \\ 0 & 0 & \xi^2 - t & -1 & 0 & \dots & 0 \\ \dots & & & & \\ 0 & 0 & 0 & 0 & 0 & \dots & -1 \\ 0 & 0 & 0 & 0 & 0 & \dots & \xi^{n-2} - t \end{vmatrix} \\ & = (-1)^{n+1} \begin{vmatrix} -1 & 0 & 0 & 0 & \dots & 0 \\ \xi - t & -1 & 0 & 0 & \dots & 0 \\ \xi - t & -1 & 0 & 0 & \dots & 0 \\ 0 & \xi^2 - t & -1 & 0 & \dots & 0 \\ \dots & & & & \\ 0 & 0 & 0 & 0 & \dots & -1 \end{vmatrix} \\ & + (\xi^{n-1} - t)(1 - t)(\xi - t) \cdots (\xi^{n-2} - t) \\ & = (-1)^{n+1}(-1)^{n-2} + (-1)^n(t - 1)(t - \xi) \cdots (t - \xi^{n-1}). \end{split}$$

Since  $\xi$  is a primitive *n*-th root of unity, the numbers  $1, \xi, \ldots, \xi^{n-1}$  are precisely all the *n*-th roots of unity, so we have

$$\chi(t) = -1 + (-1)^n (t^n - 1) = \begin{cases} -t^n & \text{if } n \text{ is odd,} \\ t^n - 2 & \text{if } n \text{ is even} \end{cases}$$

Now recall that, by Cayley–Hamilton Theorem,  $\chi(A) = 0$ , which immediately allows us to find  $A^n$ .

**Answer**:  $A^n = 0$  if n is odd and  $A^n = 2I$  if n is even.

**Remark**. An alternative solution is to observe that A = X - Y where

$$X = \begin{pmatrix} 1 & 0 & 0 & \dots & 0 & 0 \\ 0 & \xi & 0 & \dots & 0 & 0 \\ 0 & 0 & \xi^2 & \dots & 0 & 0 \\ \dots & & & & \\ 0 & 0 & 0 & \dots & \xi^{n-2} & 0 \\ 0 & 0 & 0 & \dots & 0 & \xi^{n-1} \end{pmatrix} \text{ and } Y = \begin{pmatrix} 0 & 1 & 0 & 0 & \dots & 0 \\ 0 & 0 & 1 & 0 & \dots & 0 \\ 0 & 0 & 0 & 1 & \dots & 0 \\ \dots & & & & & \\ 0 & 0 & 0 & 0 & \dots & 1 \\ 1 & 0 & 0 & 0 & \dots & 0 \end{pmatrix}.$$

One immediately verifies that  $YX = \xi XY$ , hence we can apply the Quantum Binomial Formula to compute:

$$A^{n} = (X - Y)^{n} = X^{n} + \sum_{k=1}^{n-1} {n \brack k}_{\xi} (-1)^{k} X^{n-k} Y^{k} + (-1)^{n} Y^{n}$$

where  $\begin{bmatrix} n \\ k \end{bmatrix}_{\xi}$  are the Gaussian binomial coefficients. Since  $\xi$  is a primitive *n*-th root of unity, all these coefficients vanish for  $k = 1, \ldots, n-1$ . It remains to observe that  $X^n = Y^n = I$ .

**3**. Harry Potter and Voldemort are playing a game by filling the squares of a chess board with integers as follows. First Voldemort fills the dark squares and then Harry must fill the light squares so that the resulting  $8 \times 8$  matrix has rank r. For what values of r does Harry have a winning strategy?

**Solution**: It is easier to see what is going on if we permute the rows and columns of the chessboard so it looks like

1	1	1	1	0	0	0	0	
1	1	1	1	0	0	0	0	
1	1	1	1	0	0	0	0	
1	1	1	1	0	0	0	0	
0	0	0	0	1	1	1	1	
0	0	0	0	1	1	1	1	
0	0	0	0	1	1	1	1	
0	0	0	0	1	1	1	1	

where 0 stands for a light square and 1 for a dark square. Thus, during the game, Voldemort replaces the 1's by integers of his choice and Harry does the same with the 0's.

Now it is clear that, for r < 4, Harry does not have a winning strategy because Voldemort can create a nonzero minor of size 4. Let us show that, for any  $r \ge 4$ , Harry does have a winning strategy. For a given  $r \ge 4$  and any  $4 \times 4$  matrices A and B, Harry must find  $4 \times 4$  matrices X and Y such that  $\operatorname{rk} \begin{bmatrix} A & X \\ Y & B \end{bmatrix} = r$ . One solution is X = I and Y = C + BA where C is any  $4 \times 4$  matrix of rank r - 4. Indeed,

$$\operatorname{rk} \begin{bmatrix} A & I \\ C + BA & B \end{bmatrix} = \operatorname{rk} \left( \begin{bmatrix} A & I \\ C + BA & B \end{bmatrix} \begin{bmatrix} I & 0 \\ -A & I \end{bmatrix} \right) = \operatorname{rk} \begin{bmatrix} 0 & I \\ C & B \end{bmatrix}$$
$$= \operatorname{rk} \left( \begin{bmatrix} I & 0 \\ -B & I \end{bmatrix} \begin{bmatrix} 0 & I \\ C & B \end{bmatrix} \right) = \operatorname{rk} \begin{bmatrix} 0 & I \\ C & 0 \end{bmatrix} = 4 + \operatorname{rk} C$$
$$= r.$$

We have used the fact that multiplication by a nonsingular matrix on the left or on the right does not change the rank.

**Answer**: Harry has a winning strategy if and only if  $4 \le r \le 8$ .

4. Let f(z) be a polynomial function with complex coefficients. Prove that if f maps the set of all complex roots of unity to itself, then f is a monomial.

**Solution**: Write  $f(z) = \sum_{k} a_k z^k$ . For any root of unity  $\xi$ , the value  $f(\xi)$  is also a root of unity, so  $f(\xi)\overline{f(\xi)} = 1$ . We have  $\overline{f(\xi)} = \overline{f(\xi)}$  where  $\overline{f(z)} = \sum_{k} \overline{a}_k z^k$ 

is the conjugate polynomial. Also,  $\bar{\xi} = \xi^{-1}$ . Consider the Laurent polynomial  $g(z) = f(z)\bar{f}(z^{-1}) - 1$ . By the above argument, all complex roots of unity are roots of g(z). It follows that g(z) is identically zero. Hence  $f(z)\bar{f}(z^{-1}) = 1$ .

For a Laurent polynomial  $h(z) = \sum_{k=s}^{t} c_k z^k$  with  $c_s \neq 0$  and  $c_t \neq 0$ , let  $\operatorname{HT}(h) = c_t z^t$  (highest term) and  $\operatorname{LT}(h) = c_s z^s$  (lowest term). Then we have  $\operatorname{HT}(h_1h_2) = \operatorname{HT}(h_1)\operatorname{HT}(h_2)$  and  $\operatorname{LT}(h_1h_2) = \operatorname{LT}(h_1)\operatorname{LT}(h_2)$ . Applying this to the equation  $f(z)\overline{f}(z^{-1}) = 1$ , we conclude that  $\operatorname{HT}(f) = \operatorname{LT}(f)$ , that is, f(z) is a monomial.

**5**. Let p be a prime. For any integers  $a_0, a_1, \ldots, a_p$ , prove that

$$\sum a_{i_1} \cdots a_{i_p} \equiv a_1^p \pmod{p}$$

where the summation is over all  $i_1, \ldots, i_p \ge 0$  such that  $i_1 + \cdots + i_p = p$ .

**Solution**: It will be convenient to work in the field  $\mathbb{Z}_p$  of integers modulo p. Let t be an indeterminate and let  $f = \sum_{k=0}^{p} a_k t^k \in \mathbb{Z}_p[t]$ . Then the sum  $\sum a_{i_1} \cdots a_{i_p}$  equals the coefficient of  $t^p$  in the polynomial  $f^p$ . But we have  $(a+b)^p = a^p + b^p$  in characteristic p, hence, by induction,  $f^p = \sum_{k=0}^{p} a_k^p t^{kp}$ , so the coefficient of  $t^p$  equals  $a_1^p$ .

**6.** Let G be a finite simple group. Suppose G contains a subgroup of prime index p. Prove that p is the largest prime divisor of |G| and, moreover,  $p^2$  does not divide |G|.

**Solution**: Let  $H \subset G$  be a subgroup of index p and let X = G/H be the set of left cosets of G with respect to H. Then G acts on X by left translations. Since |X| = p, this action yields a homomorphism from G to  $S_p$ , the symmetric group on p symbols. Since the action of G on X is nontrivial (in fact, transitive), the kernel of the homomorphism  $G \to S_p$  cannot be the entire G, hence it must be trivial by simplicity. Therefore, G is isomorphic to a subgroup of  $S_p$ . Hence |G| is a divisor of p!. The result follows.